

Topic 1 Analytic Combinatorics

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Enumerative Combinatorics

Start with a sequence

$$(f_n) = f_0, f_1, f_2, \dots$$

The terms of the sequence could

- count objects in a combinatorial class
- capture the **probability** that an event occurs
- track the **runtime** of an algorithm

Goal: capture something interesting about the sequence

Exact Formulas

A THEOREM ON TREES.

By Prof. CAYLEY.

THE number of trees which can be formed with n+1 given knots α , β , γ , ... is $=(n+1)^{n-1}$; for instance n=3, the number of trees with the 4 given knots α , β , γ , δ is $4^2=16$

A. Cayley. A Theorem on Trees. Quart. J. Pure Appl. Math. Vol 23, 376-378, 1889.

It's unreasonable to expect this to always occur — not all combinatorial sequences have *simple* formulas, and even if they do they can be hard to prove!

Algorithms

25 loops, best of 3: 10.1 ms per loop

Gathering data can be useful for studying sequences, and conjecturing formulas, but doesn't fully *capture behaviour*.

Asymptotics

Instead of exact enumeration, focus on large-scale behaviour by approximating f_n for large n.

partitions of
$$n \sim \frac{1}{4n\sqrt{3}} \exp\left(\pi\sqrt{\frac{2n}{3}}\right)$$

average quicksort cost on permutation of size n $\sim 2n \log n$

unlabelled graphs on
$$n$$
 nodes $\sim \frac{2^{\binom{2n}{n}}}{n!}$

The generating function (GF) of f_n is

$$F(z) = \sum_{n \ge 0} f_n z^n$$

Algebraic / differential / functional equations for F form a data structure for f_n

 $f_n = \#$ walks starting at $\mathbf{0}$ on n steps in $\mathbf{4}$

$$F(z) = \frac{1}{1 - 4z}$$

The generating function (GF) of f_n is

$$F(z) = \sum_{n \ge 0} f_n z^n$$

Algebraic / differential / functional equations for F form a data structure for f_n

 $f_n = \#$ walks starting at $\mathbf{0}$ on $\stackrel{\uparrow}{\longleftrightarrow}$ staying in halfspace

$$F(z) = \frac{4z + \sqrt{1 - 4z} - 1}{2z(1 - 4z)}$$

The generating function (GF) of f_n is

$$F(z) = \sum_{n>0} f_n z^n$$

Algebraic / differential / functional equations for F form a data structure for f_n

 $f_n = \#$ walks starting at $\mathbf{0}$ on \longleftrightarrow staying in $3\pi/4$ wedge

$$27z^{7}(4z-1)^{2}F(z)^{8} + \dots + (16z^{2} - 12z + 1)F(z) - 1 = 0$$

The generating function (GF) of f_n is

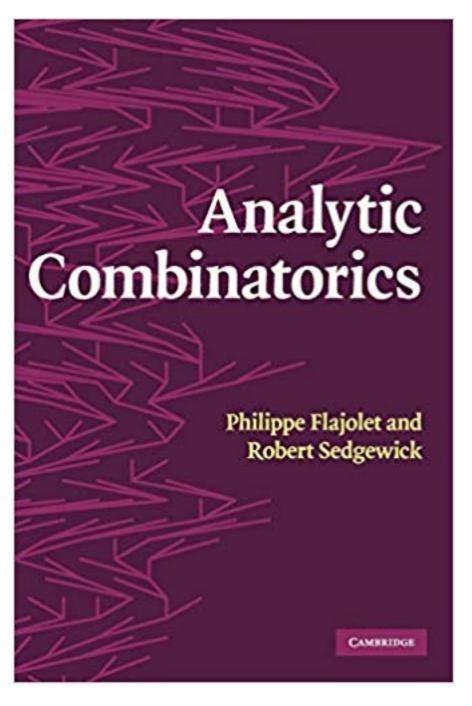
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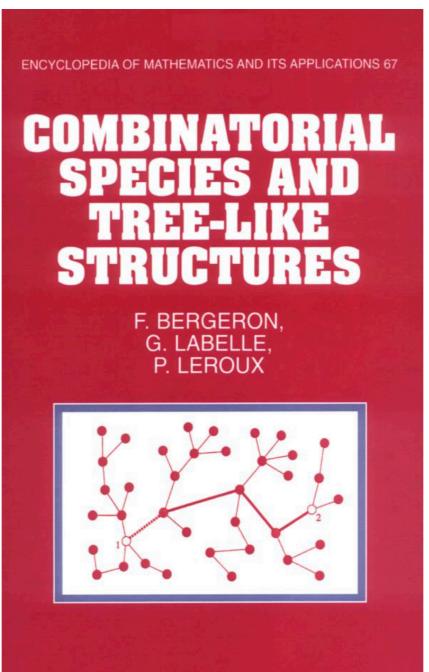
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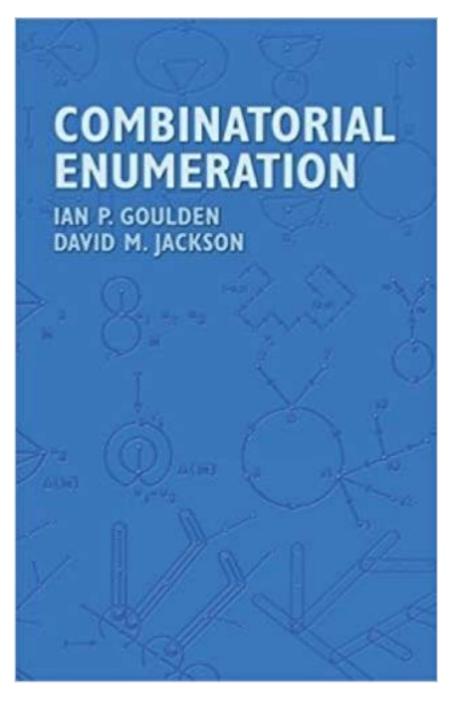
 $f_n = \#$ walks starting at $\mathbf{0}$ on \longleftrightarrow staying in quadrant

$$z^{2}(4z-1)(4z+1)F'''(z) + 2z(4z+1)(16z-3)F''(z) + 2(112z^{2} + 14z - 3)F'(z) + 4(16z+3)F(z) = 0$$

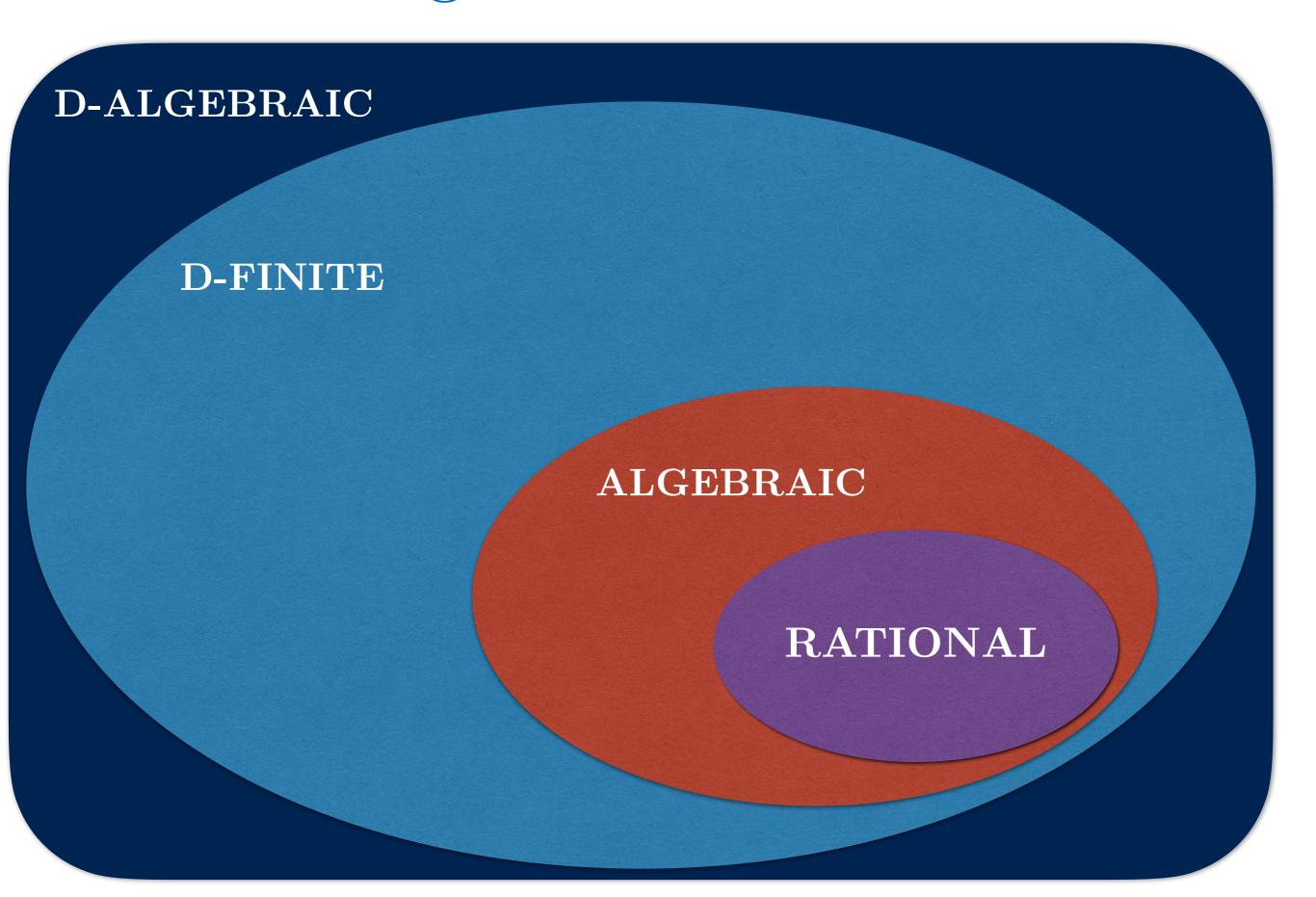
Combinatorial definitions often automatically translate to GF equations







Generating Function Classes



Analytic Combinatorics

We assume our GF is **analytic** at the origin



$$\sum_{n\geq 0} 2^n z^n$$

$$\sum_{n\geq 0} \frac{z^n}{n}$$

$$\sum_{n\geq 0} n! \, z^n$$

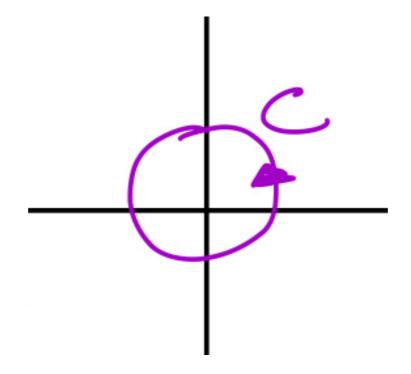
Analytic combinatorics derives asymptotics of sequences from the behaviour of their GFs

Fact 1: Cauchy Integral Formula

If $F(z) = \sum_{n>0} f_n z^n$ is analytic at the origin then

$$f_n = \frac{1}{2\pi i} \int_{\mathcal{C}} \frac{F(z)}{z^{n+1}} dz$$

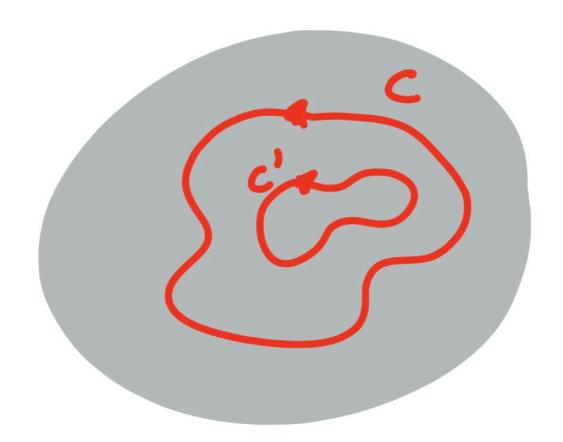
where \mathcal{C} is a sufficiently small circle around the origin.



Fact 2: Deforming Curves of Integration

If C and C' are simple closed curves and C can be deformed to C' in an open set where f(z) is analytic then

$$\int_{\mathcal{C}} f(z)dz = \int_{\mathcal{C}'} f(z)dz$$



Fact 3: Residues



Assume

- P(z) and Q(z) analytic at $z = \rho$
- \mathcal{C} is any sufficiently small circle around $z = \rho$

Then

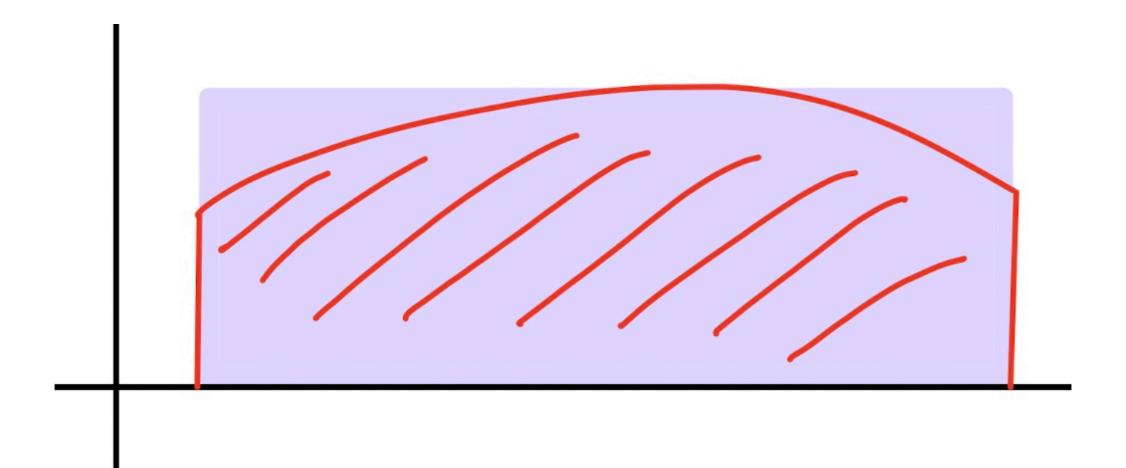
$$\int_{\mathcal{C}} \frac{P(z)}{Q(z)} dz = \operatorname{Res}_{z=\rho} \left(\frac{P(z)}{Q(z)} \right)$$

where the **residue** is an explicit and computable expression involving the derivatives of P and Q at $z = \rho$

Fact 4: Max Modulus Bound

If f(z) continuous on C then

$$\left| \int_{\mathcal{C}} f(z) dz \right| \le \operatorname{length}(\mathcal{C}) \times \max_{z \in \mathcal{C}} |f(z)|$$



Alternating Permutations

An alternating permutation is a permutation π of odd length such that $\pi_1 > \pi_2 < \pi_3 > \cdots$

The alternating permutations of length three: 213 and 312

$$A(z) = \sum_{k>0} \frac{a_{2k+1}}{(2k+1)!} z^{2k+1}$$

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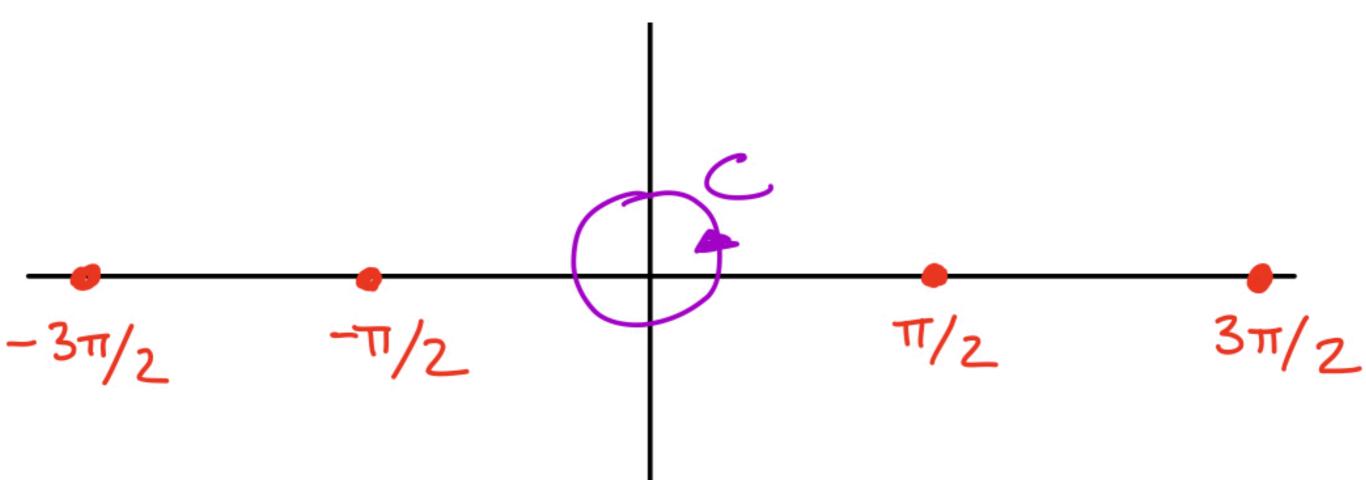
$$A(z) = \sum_{k>0} \frac{a_{2k+1}}{(2k+1)!} z^{2k+1} = \tan z$$

ANALYSE MATHÉMATIQUE. — Développements de séc x et de tang x. Note de M. D. André, présentée par M. Hermite.

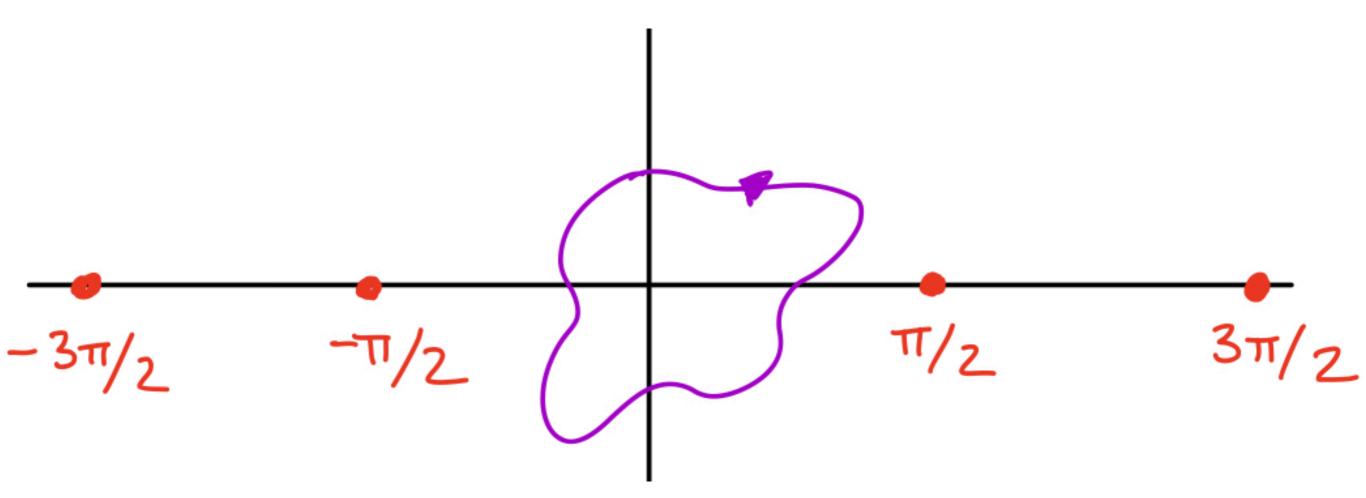
« On n'a point donné jusqu'à présent, du moins à ma connaissance, de développement, suivant les puissances de x, soit de tang x, soit de séc x, où les coefficients aient une définition simple, nette, indépendante de tout autre développement. L'objet de la présente Note est de combler cette lacune.

Comptes rendus de l'Académie des sciences, 1879

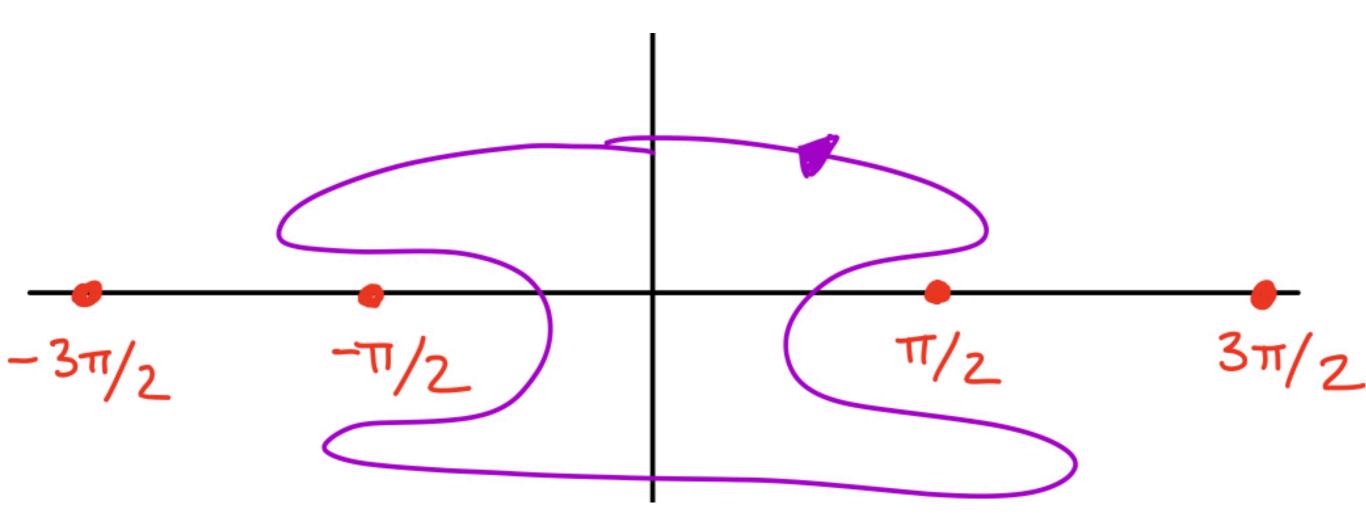
$$\frac{a_n}{n!} = [z^n] \tan z = \frac{1}{2\pi i} \int_{\mathcal{C}} \frac{\tan z}{z^{n+1}}$$



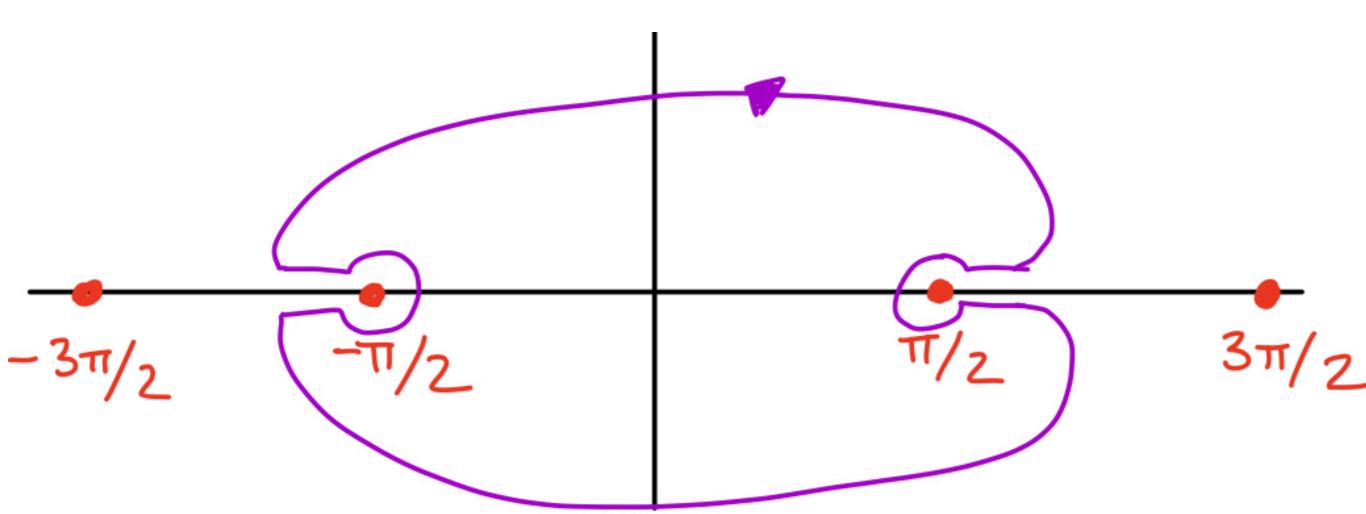
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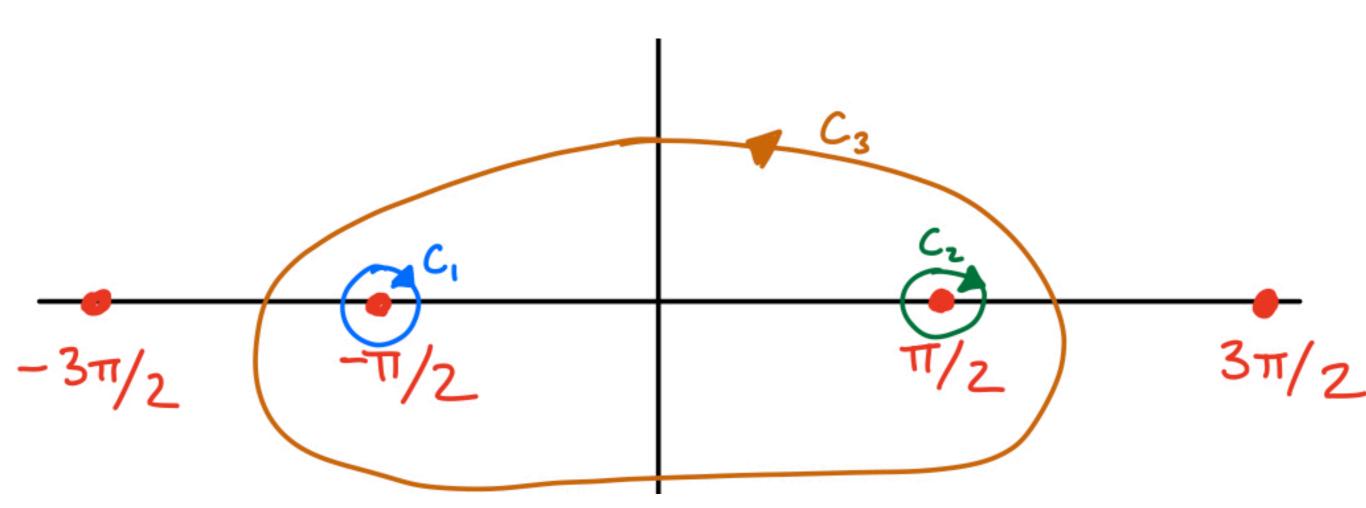
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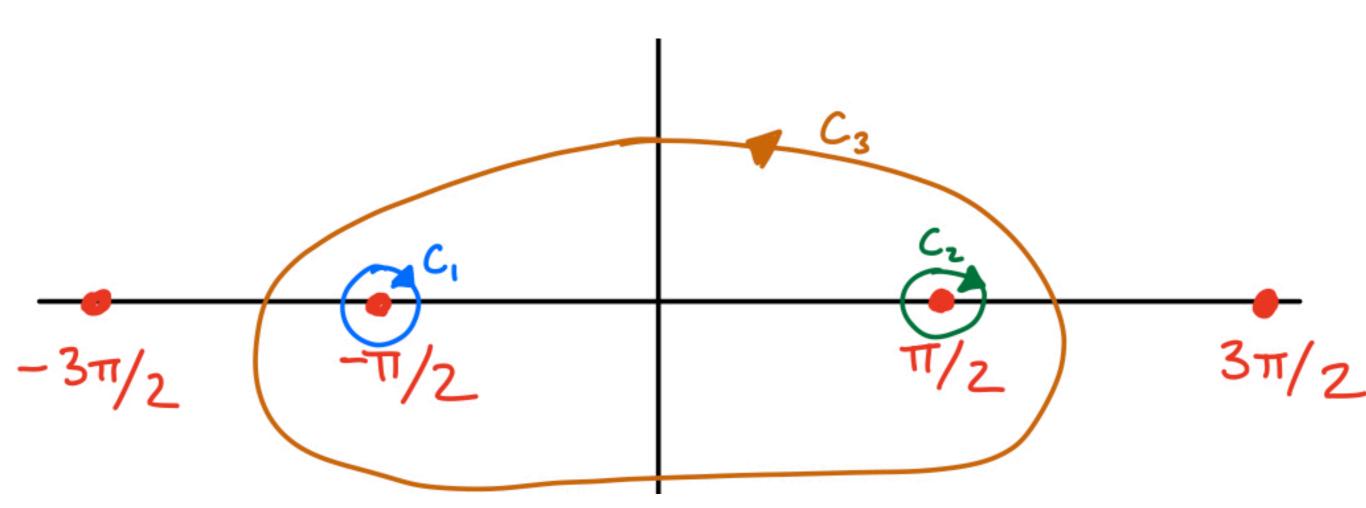
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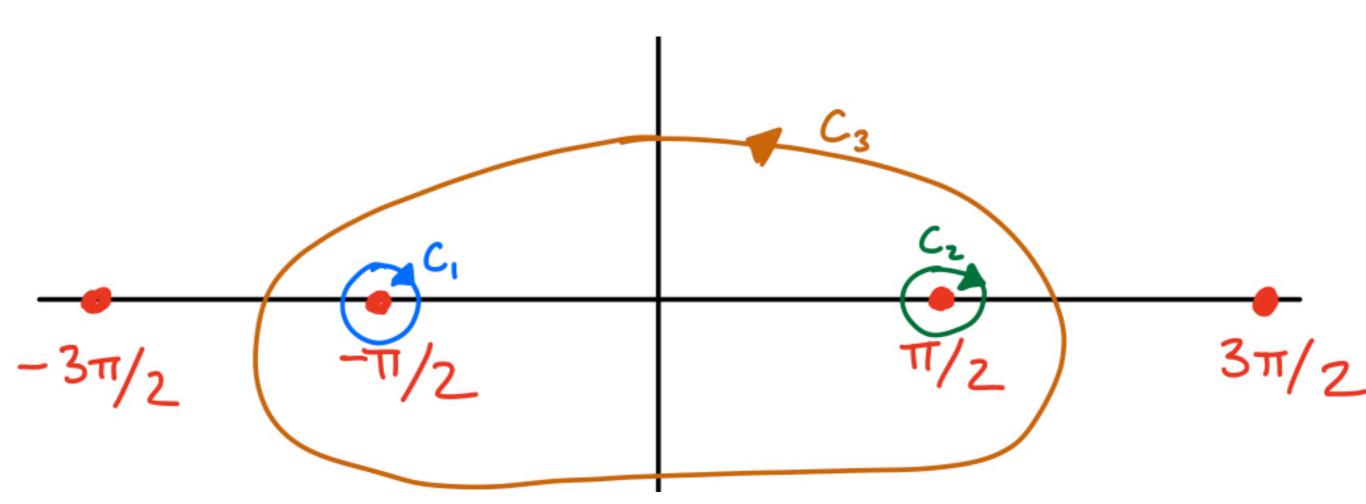
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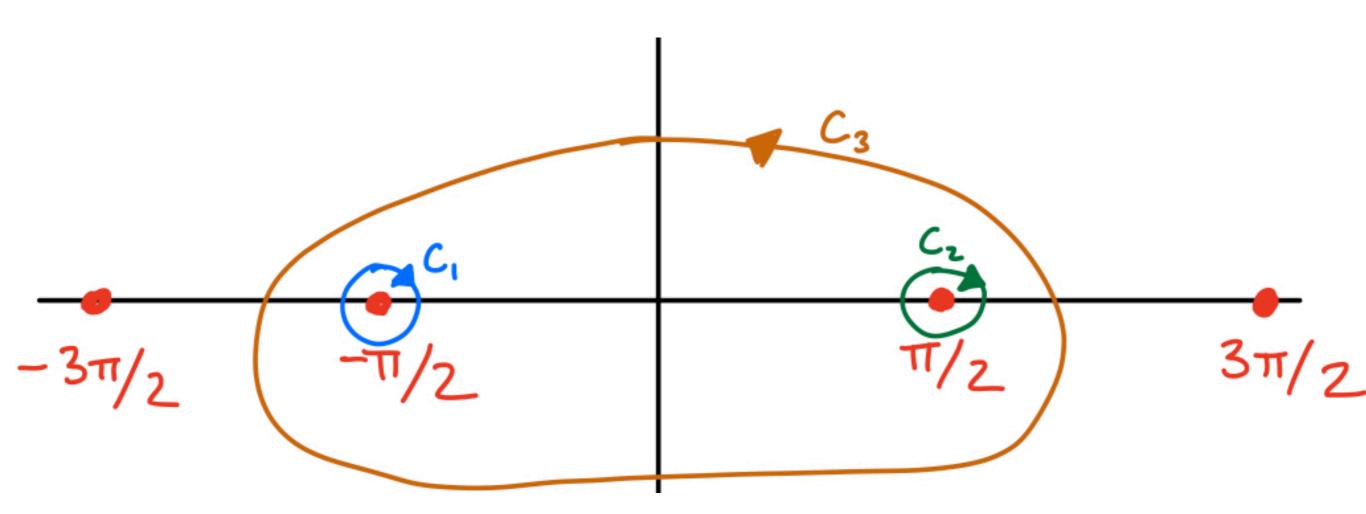
$$\frac{a_n}{n!} = \frac{1}{2\pi i} \int_{\mathcal{C}_1} \frac{\tan z}{z^{n+1}} dz + \frac{1}{2\pi i} \int_{\mathcal{C}_2} \frac{\tan z}{z^{n+1}} dz + \frac{1}{2\pi i} \int_{|z|=\pi} \frac{\tan z}{z^{n+1}} dz$$



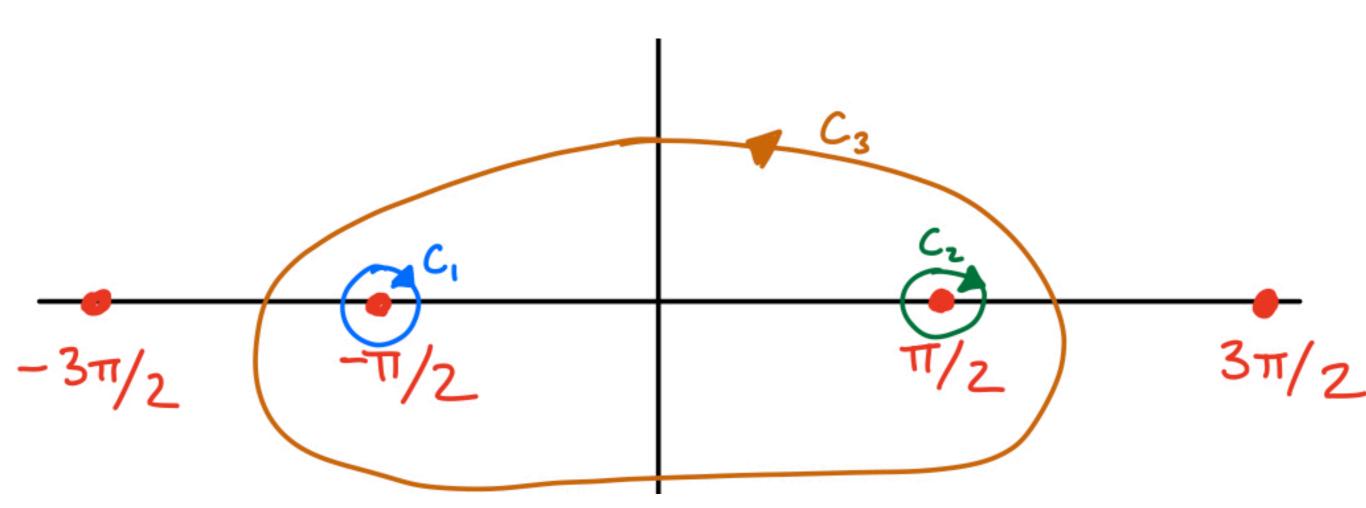
$$\frac{a_n}{n!} = \underset{z=-\pi/2}{\text{Res}} \left(\frac{\tan z}{z^{n+1}}\right) + \underset{z=\pi/2}{\text{Res}} \left(\frac{\tan z}{z^{n+1}}\right) + O\left(\left(\frac{1}{\pi}\right)^n\right)$$



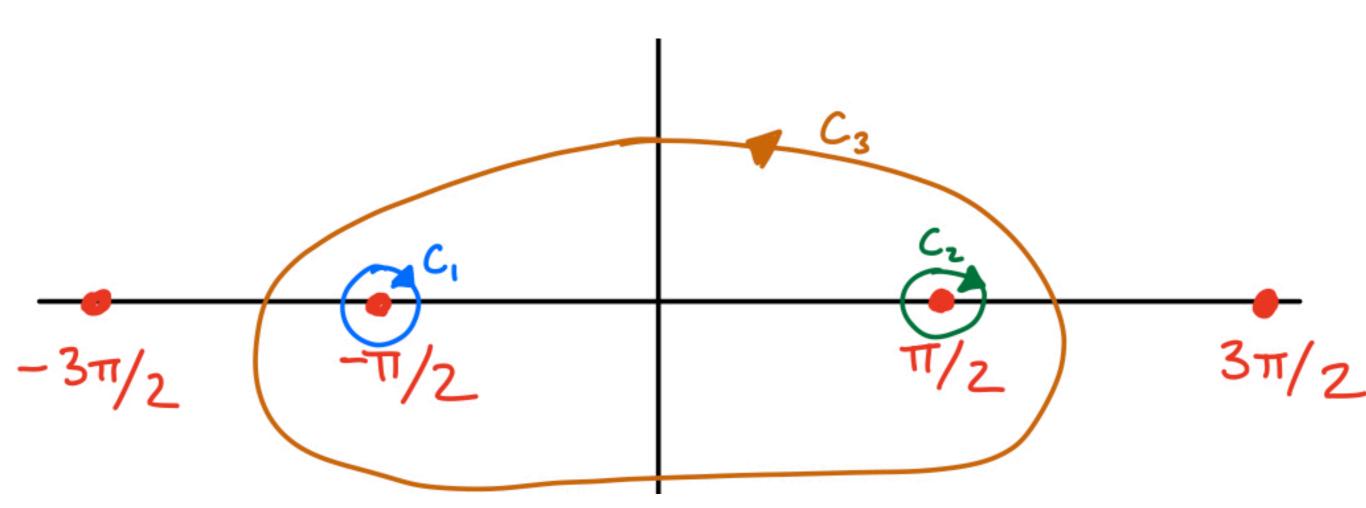
$$\frac{a_n}{n!} = \left(\frac{-2}{\pi}\right)^{n+1} \operatorname{Res}_{z=-\pi/2} \left(\frac{\sin z}{\cos z}\right) + \left(\frac{2}{\pi}\right)^{n+1} \operatorname{Res}_{z=\pi/2} \left(\frac{\sin z}{\cos z}\right) + O\left(\left(\frac{1}{\pi}\right)^n\right)$$



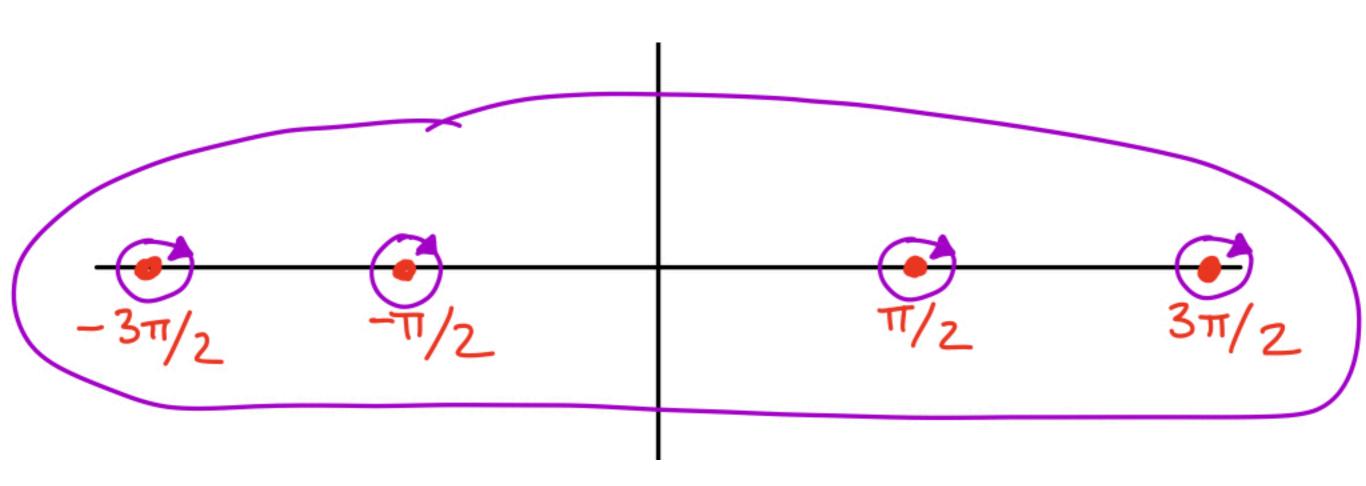
$$\frac{a_n}{n!} = \left(\frac{-2}{\pi}\right)^{n+1} + \left(\frac{2}{\pi}\right)^{n+1} + O\left(\left(\frac{1}{\pi}\right)^n\right)$$



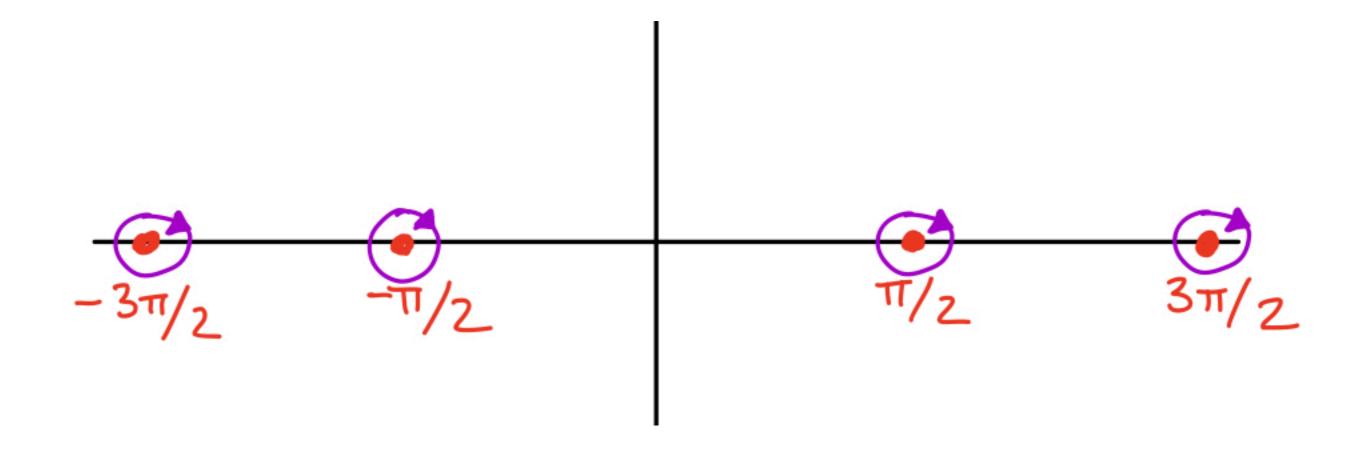
$$\frac{a_n}{n!} = 2\left(\frac{2}{\pi}\right)^{n+1} + O\left(\left(\frac{1}{\pi}\right)^n\right) \qquad (n \text{ odd})$$



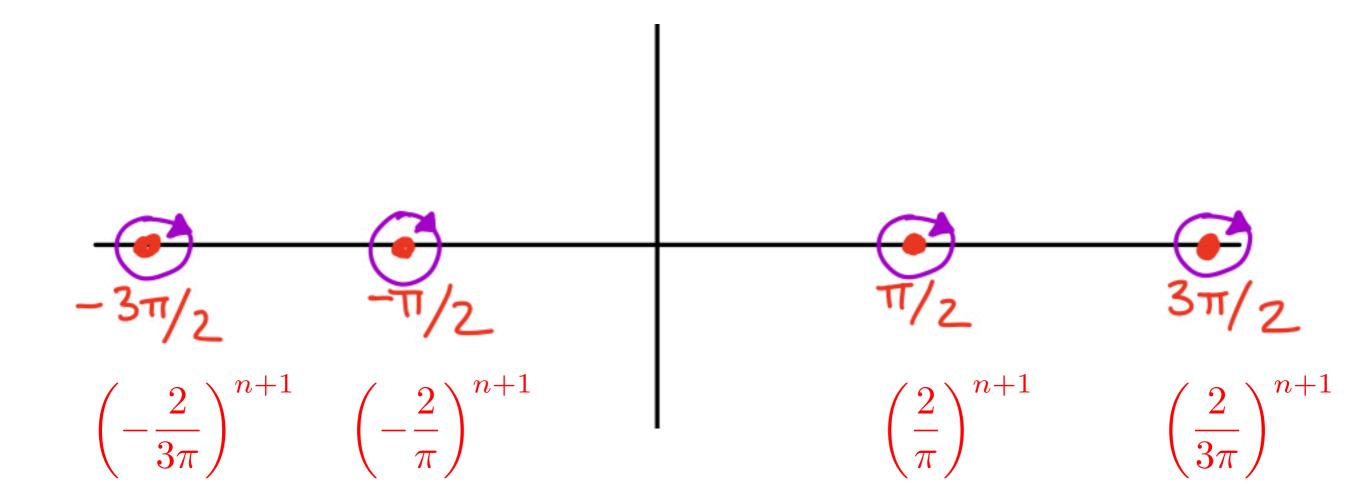
$$\frac{a_n}{n!} = 2\left(\frac{2}{\pi}\right)^{n+1} + 2\left(\frac{2}{3\pi}\right)^{n+1} + O\left(\left(\frac{2}{5\pi}\right)^n\right) \qquad (n \text{ odd})$$



$$\frac{a_n}{n!} = 2\left(\frac{2}{\pi}\right)^{n+1} \sum_{k>0} \frac{1}{(2k+1)^{n+1}}$$
 (n odd)



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Main Takeaways

- Each **singularity** gives contribution
- Those singularities **closest to the origin** affect dominant asymptotics
- The contributions of each can be determined by a **local analysis** of the generating function

There are many known formulas for different types of singularities

$$F(z) \sim (1-z)^{\alpha} \left(\log \frac{1}{1-z}\right)^{\beta} \implies f_n \sim \frac{n^{-\alpha-1}}{\Gamma(-\alpha)} (\log n)^{\beta}$$

Topic 2 Diagonals and Smooth ACSV

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Diagonals

Start with a multivariate series

$$F(\mathbf{z}) = \sum_{(i_1, \dots, i_d) \in \mathbb{N}^d} f_{i_1, \dots, i_d} z_1^{i_1} \cdots z_d^{i_d} = \sum_{\mathbf{i} \in \mathbb{N}^d} f_{\mathbf{i}} \mathbf{z}^{\mathbf{i}}$$

Diagonals

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The **r-diagonal** consists of the coefficients $(f_{n\mathbf{r}}) = f_{\mathbf{0}}, f_{\mathbf{r}}, f_{2\mathbf{r}}, \dots$ Note the coefficient $f_{n\mathbf{r}}$ is defined only if $n\mathbf{r} \in \mathbb{N}^d$

(1,1) – Diagonal (Main Diagonal)

$$F(x,y) = \frac{1}{1-x-y}$$

$$= 1 + x + y + 2xy + x^2 + y^2 + 3x^2y + 3xy^2 + y^3 + 6x^2y^2 + \cdots$$

Diagonals

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$$(2,1) - \textbf{Diagonal}$$

$$F(x,y) = \frac{1}{1-x-y}$$

$$= 1 + x + y + 2xy + x^2 + y^2 + 3x^2y + \dots + 15x^4y^2 + \dots$$

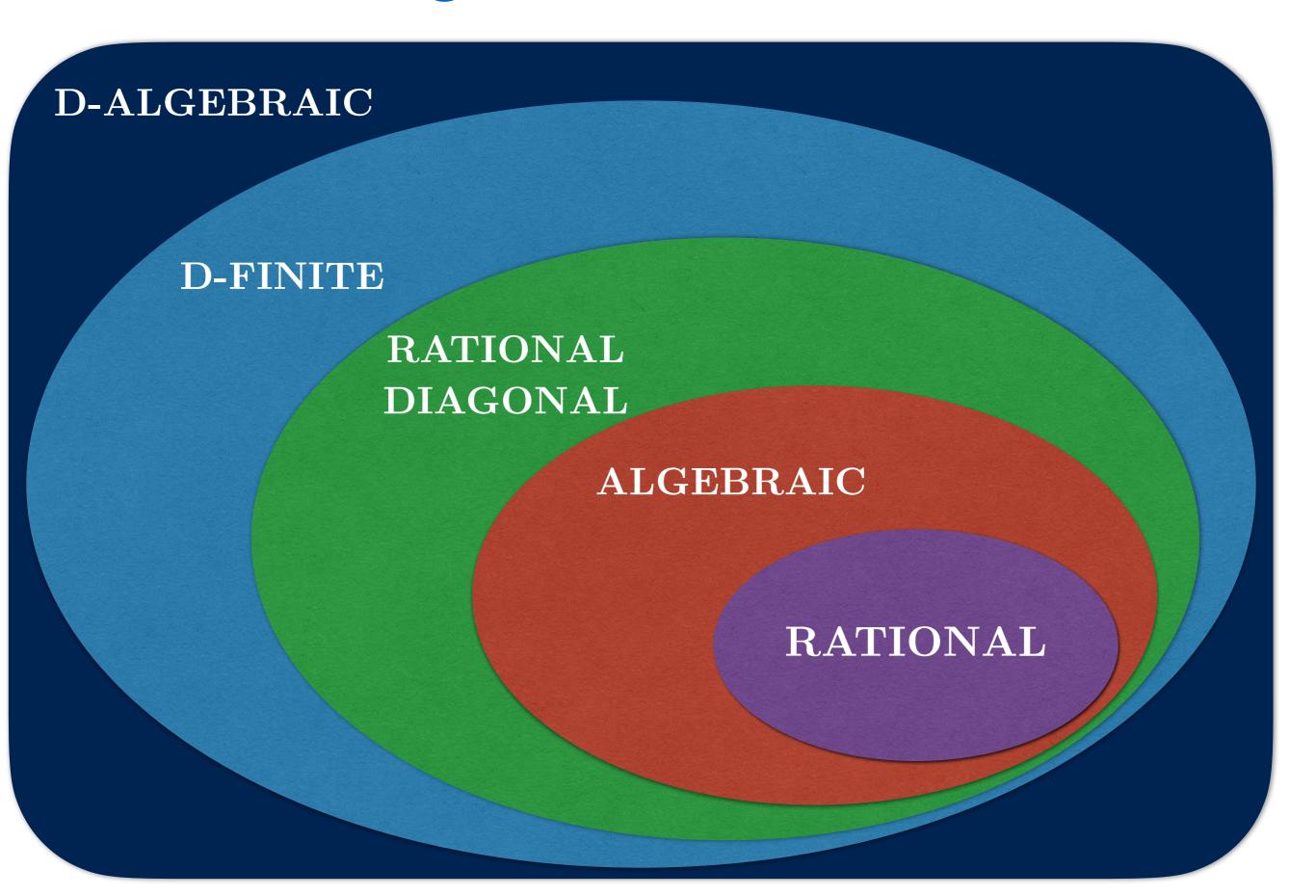
Why Diagonals?

- Data structures for interesting univariate sequences
- Uniform asymptotics over *most* directions (tomorrow)
- Yield combinatorial limit theorems (tomorrow)

We focus on **rational** (or **meromorphic**) diagonals

• Diagonal of an algebraic function in d variables is the diagonal of a rational function in 2d variables

Generating Function Classes



Analytic Combinatorics in Several Variables

Assume

$$F(\mathbf{z}) = \frac{G(\mathbf{z})}{H(\mathbf{z})} = \sum_{\mathbf{i} \in \mathbb{N}^d} f_{\mathbf{i}} \mathbf{z}^{\mathbf{i}}$$

coprime ratio

converges in a neighbourhood of the origin.

The singularities of $F(\mathbf{z})$ are given by $\mathcal{V} := {\mathbf{z} : H(\mathbf{z}) = 0}$. Singularities **closest** to the origin are called **minimal points**.

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 \mathbf{w} minimal if and only if $H(\mathbf{w}) = 0$ and there is no \mathbf{z} with

$$H(\mathbf{z}) = 0$$
 and $|z_j| < |w_j|$ for all j

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Multivariate Cauchy Integral Formula

$$f_{n\mathbf{r}} = \frac{1}{(2\pi i)^d} \int_{\mathcal{C}} F(\mathbf{z}) \frac{d\mathbf{z}}{\mathbf{z}^{n\mathbf{r}+1}} \qquad n\mathbf{r} \in \mathbb{N}^d$$

where \mathcal{C} is a product of circles $|z_i| = \varepsilon$

Difficulties of ACSV

One variable rational (or meromorphic) functions

- Find finite set of singularities closest to the origin
- Add their asymptotic contributions

In more than one variable

- Set of minimal points is infinite
- Singular set can have nontrivial geometry (self-intersections)
- Can deform domain of integration around singular set!

Smooth ACSV

Simplest case: Denominator H and its partial derivatives don't simultaneously vanish.

Then **critical points** are defined by

$$H = 0,$$
 $r_j z_1 H_{z_1} = r_1 z_j H_{z_j}$ $(2 \le j \le d)$

partial derivative

Smooth ACSV

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Critical points: Asymptotic approximations can be made Minimal points: Cauchy integral can be deformed close to

The asymptotic contribution of a minimal critical point \mathbf{w} depends on an explicit matrix $\mathcal{M} = \mathcal{M}_{\mathbf{w}}$ built from partial derivatives of H

Surgery ACSV Theorem (Pemantle Wilson 2003)

Suppose that

$$H = 0,$$
 $r_j z_1 H_{z_1} = r_1 z_j H_{z_j}$ $(2 \le j \le d)$

admits a minimal solution $\mathbf{w} \in \mathbb{C}^d_*$. If

- no other singularity has the same coordinate-wise modulus as **w**
- $H_{z_d}(\mathbf{w})$ and $\det \mathcal{M}$ are non-zero,

then

$$\left[\mathbf{z}^{n\mathbf{r}}\right] \frac{G(\mathbf{z})}{H(\mathbf{z})} = \mathbf{w}^{-n\mathbf{r}} (nr_d)^{(1-d)/2} (2\pi)^{(1-d)/2} \det(\mathcal{M})^{-1/2} \left(\frac{-G(\mathbf{w})}{w_d H_{z_d}(\mathbf{w})} + O\left(\frac{1}{n}\right)\right)$$

If there are a finite number of singularities with the same coordinate-wise modulus as \mathbf{w} , all satisfying these conditions, then we can add their asymptotic contributions.

If $H_{z_d}(\mathbf{w}) \neq 0$ then we can write $z_d = g(\widehat{\mathbf{z}})$ near \mathbf{w}

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 \mathcal{M} is the $(d-1)\times(d-1)$ Hessian matrix at $\theta=\mathbf{0}$ of

$$\phi(\theta) = \log g\left(w_1 e^{i\theta_1}, \dots, w_{d-1} e^{i\theta_{d-1}}\right)$$

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The Chain Rule implies

$$\mathcal{M}_{i,j} = \begin{cases} V_i V_j + U_{i,j} - V_j U_{i,d} - V_i U_{j,d} + V_i V_j U_{d,d} & : i \neq j \\ V_i + V_i^2 + U_{i,i} - 2V_i U_{i,d} + V_i^2 U_{d,d} & : i = j \end{cases}$$

where
$$U_{i,j} = \frac{w_i w_j H_{z_i z_j}(\mathbf{w})}{w_d H_{z_d}(\mathbf{w})}$$
 and $V_i = r_i/r_d$

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```
def getHes(H,R,vars,CP):
    dd = len(vars)
    V = zero_vector(SR,dd)
    U = matrix(SR,dd)
    M = matrix(SR,dd-1)

for j in range(dd):
        V[j] = R[j]/R[-1]
        for i in range(dd):
              U[i,j] = vars[i]*vars[j]*diff(H,vars[i],vars[j])/vars[-1]/diff(H,vars[-1])

for i in range(dd-1):
    for j in range(dd-1):
        M[i,j] = V[i]*V[j] + U[i,j] - V[j]*U[i,-1] - V[i]*U[j,-1] + V[i]*V[j]*U[-1,-1]
        if i == j: M[i,j] = M[i,j] + V[i]
    return M.subs(CP)
```

A First Example

$$F(x,y) = \frac{1}{1 - x - y} = \sum_{i,j>0} {i+j \choose i} x^{i} y^{j}$$

Critical Point Equations in Direction r = (1, 1)

$$1 - x - y = 0 \qquad \qquad -x = -y$$

Unique Minimal Critical Point

$$(x_*, y_*) = \left(\frac{1}{2}, \frac{1}{2}\right)$$

Hessian

$$g(x) = 1 - x$$
 so $\phi(\theta) = \log\left(1 - \frac{1}{2}e^{i\theta}\right)$ and $\mathcal{M} = \phi''(0) = 2$

A First Example

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Asymptotics

$$\binom{2n}{n} = [x^n y^n] F(x, y) = \frac{4^n}{\sqrt{\pi n}} (1 + O(n^{-1}))$$

A First Example

$$F(x,y) = \frac{1}{1 - x - y} = \sum_{i,j>0} {i+j \choose i} x^{i} y^{j}$$

Critical Point Equations in Direction $\mathbf{r} = (r, s)$

$$1 - x - y = 0 \qquad -sx = -ry$$

Unique Minimal Critical Point

$$(x_*, y_*) = \left(\frac{r}{r+s}, \frac{s}{r+s}\right)$$

Asymptotics

$$\binom{rn+sn}{rn} = [x^{rn}y^{sn}]F(x,y) = \left(\frac{r+s}{r}\right)^{rn} \left(\frac{r+s}{s}\right)^{sn} \frac{\sqrt{r+s}}{\sqrt{2rs\pi n}} \left(1 + O\left(n^{-1}\right)\right)$$

Bi-Clover Quiver (Ramgoolam, Wilson and Zahabi 2020)

The generating function for the chiral operators in the large N limit of the bi-clover quiver gauge theory is

$$F(x,y) = \frac{1}{\prod_{k>1} (1 - x^k - y^k)}$$

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$$F(x,y) = \frac{1}{\prod_{k>1} (1 - x^k - y^k)}$$

Note: F(x,y) = G(x,y)/(1-x-y) where

$$G(x,y) = \prod_{k>2} (1 - x^k - y^k)^{-1}$$

Thus,

$$[x^{rn}y^{sn}]F(x,y) = G\left(\frac{r}{r+s}, \frac{s}{r+s}\right)\left(\frac{r+s}{r}\right)^{rn}\left(\frac{r+s}{s}\right)^{sn}\frac{\sqrt{r+s}}{\sqrt{2rs\pi n}}\left(1 + O\left(n^{-1}\right)\right)$$

Suppose $\mathbf{w} = (a, b) \in \mathbb{R}^2_{>0}$ satisfies the conditions of the theorem.

Cauchy Integral Formula implies

$$f_{rn,sn} = \frac{1}{(2\pi i)^2} \int_{|x|=a} \left(\int_{|y|=b-\varepsilon} F(x,y) \frac{dx \, dy}{x^{rn+1} y^{sn+1}} \right)$$

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Cauchy Integral Formula and Max Modulus Bound imply

$$f_{rn,sn} = \frac{1}{(2\pi i)^2} \int_{|x|=a} \left(\int_{|y|=b-\varepsilon} F(x,y) \frac{dx \, dy}{x^{rn+1} y^{sn+1}} \right)$$
$$-\frac{1}{(2\pi i)^2} \int_{|x|=a} \left(\int_{|y|=b+\varepsilon} F(x,y) \frac{dx \, dy}{x^{rn+1} y^{sn+1}} \right) + \text{small error}$$

^{*} May need to localize

Suppose $\mathbf{w} = (a, b) \in \mathbb{R}^2_{>0}$ satisfies the conditions of the theorem.

Cauchy Integral Formula and Max Modulus Bound and Residues imply

$$f_{rn,sn} = \frac{1}{2\pi i} \int_{|x|=a} \left(\operatorname{Res}_{y=g(x)} \frac{F(x,y)}{y^{sn+1}} \right) \frac{dx}{x^{rn+1}} + \text{small error}$$

where y = g(x) on \mathcal{V}

Suppose $\mathbf{w} = (a, b) \in \mathbb{R}^2_{>0}$ satisfies the conditions of the theorem.

Cauchy Integral Formula and Max Modulus Bound and Residues imply

$$f_{rn,sn} = \frac{1}{2\pi i} \int_{|x|=a} \frac{G(x,g(x))}{H_y(x,g(x))} \cdot \frac{dx}{x^{rn+1}g(x)^{sn+1}} + \text{small error}$$

where y = g(x) on \mathcal{V}

Suppose $\mathbf{w} = (a, b) \in \mathbb{R}^2_{>0}$ satisfies the conditions of the theorem.

Cauchy Integral Formula and Max Modulus Bound and Residues and the change of variables $x = ae^{i\theta}$ imply

$$f_{rn,sn} = \int_{-\pi}^{\pi} A(\theta)e^{-n\phi(\theta)}d\theta + \text{small error}$$

Because w is critical this is a saddle-point integral

Critical vs Minimal Points

Minimal points make the proof relative **easy**, but make checking the conditions **difficult**.

Problem: Dealing with minimality (and points with same coordinate-wise modulus) is hard. It also considers spurious points.

$$F(x,y) = \frac{1}{(1+2y)(1-x-y)}$$
 still has critical point $(1/2, 1/2)$

But now there is a **curve** of singularities

$$\{(e^{i\theta}/2, -1/2) : \theta \in (-\pi, \pi]\}$$

with
$$|x| = |y| = 1/2$$

Critical vs Minimal Points

Minimal points make the proof relative **easy**, but make checking the conditions **difficult**.

Problem: Dealing with minimality (and points with same coordinate-wise modulus) is hard. It also considers spurious points.

Solution: Reduce importance of minimality. Only *critical points* really matter when computing asymptotics.

Key: Generically there are a finite set of critical points, encoded by algebraic equations, even though there are infinite minimal points.

Main Theorem of Smooth ACSV (Baryshnikov Pemantle 2011 / BMP 2021)

Suppose that

$$H = 0,$$
 $r_j z_1 H_{z_1} = r_1 z_j H_{z_j}$ $(2 \le j \le d)$

admits a finite number of solutions. If

- there is exactly one minimal solution, $\mathbf{w} \in \mathbb{C}^d_*$
- $H_{z_d}(\mathbf{w})$ and $\det \mathcal{M}$ are non-zero,

then

$$\left[\mathbf{z}^{n\mathbf{r}}\right] \frac{G(\mathbf{z})}{H(\mathbf{z})} = \mathbf{w}^{-n\mathbf{r}} (nr_d)^{(1-d)/2} (2\pi)^{(1-d)/2} \det(\mathcal{M})^{-1/2} \left(\frac{-G(\mathbf{w})}{w_d H_{z_d}(\mathbf{w})} + O\left(\frac{1}{n}\right)\right)$$

In other words: We can study the (hopefully finite) set of critical points and check which are minimal, ignoring everything else

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$$[x^n y^n] \frac{1}{(1+2y)(1-x-y)} \sim \frac{1}{2} \cdot \frac{4^n}{\sqrt{\pi n}}$$

Main Theorem of Smooth ACSV (Baryshnikov Pemantle 2011 / BMP 2021)

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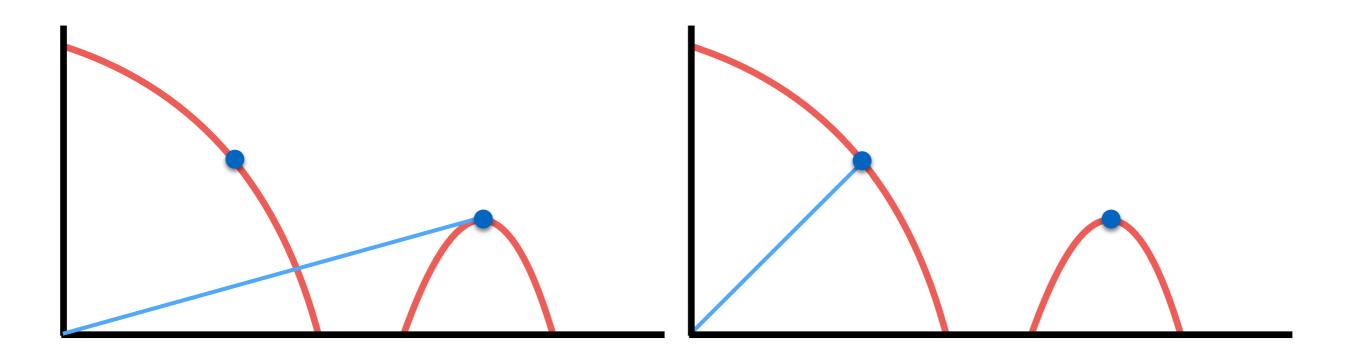
$$\left[\mathbf{z}^{n\mathbf{r}}\right] \frac{G(\mathbf{z})}{H(\mathbf{z})} = \mathbf{w}^{-n\mathbf{r}} (nr_d)^{(1-d)/2} (2\pi)^{(1-d)/2} \det(\mathcal{M})^{-1/2} \left(\frac{-G(\mathbf{w})}{w_d H_{z_d}(\mathbf{w})} + O\left(\frac{1}{n}\right)\right)$$

If there are a finite number of **critical points** with the same coordinate-wise modulus as **w**, all satisfying these conditions, then we can add their asymptotic contributions.

Help Proving Minimality

Multivariate Vivanti-Pringsheim Theorem

If $f_i \ge 0$ for all **i** then $\mathbf{w} \in \mathcal{V}$ with positive coordinates is minimal if and only if $H(t\mathbf{w}) \ne 0$ for all $t \in (0,1)$.



Help Proving Minimality

Aperiodic Expansions

If
$$H(\mathbf{z}) = 1 - \sum_{\mathbf{n} \in \mathbb{N}^d} p_{\mathbf{n}} \mathbf{z}^{\mathbf{n}}$$
 where $(p_{\mathbf{n}})$ is a sequence of

nonnegative numbers with $\operatorname{span}_{\mathbb{Z}}\{\mathbf{n} \in \mathbb{N}^d : p_{\mathbf{n}} \neq 0\} = \mathbb{Z}^d$ then every minimal point has positive real coordinates



$$\frac{1}{1-x-y}$$

$$\frac{1}{2 - e^{x+y}}$$



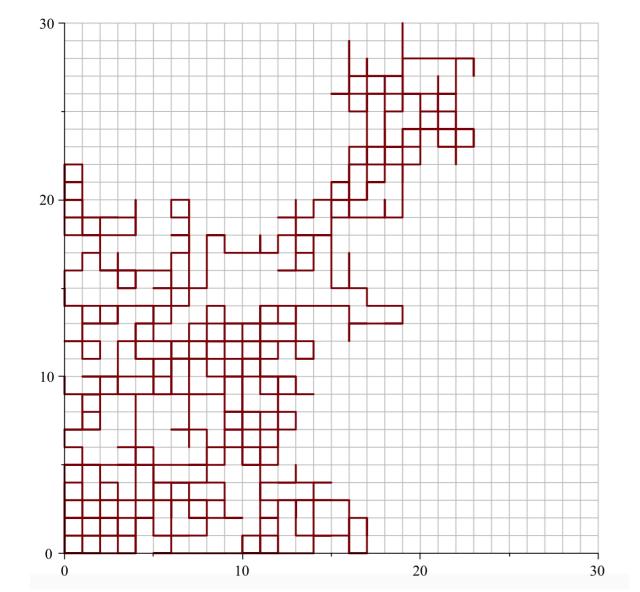
$$\frac{1}{1 - t(x+y)}$$

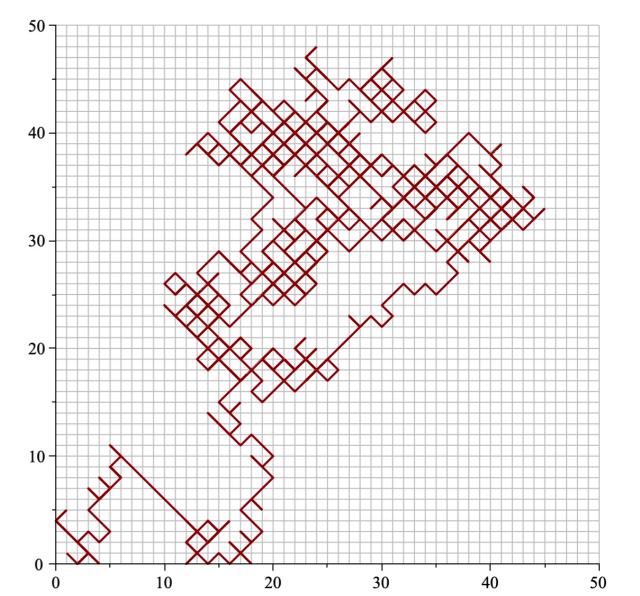
$$\frac{1}{1-x+y}$$

Walks in an Orthant

Uniform diagonal expression for walk models in \mathbb{N}^d whose step sets $\mathcal{S} \subset \{\pm 1, 0\}^d$ are symmetric over every axis.

$$[(z_1 \cdots z_d t)^n] \frac{(1+z_1) \cdots (1+z_d)}{1-t(z_1 \cdots z_d) S(\mathbf{z})}, \qquad S(\mathbf{z}) = \sum_{\mathbf{i} \in \mathcal{S}} \mathbf{z}^{\mathbf{i}}$$



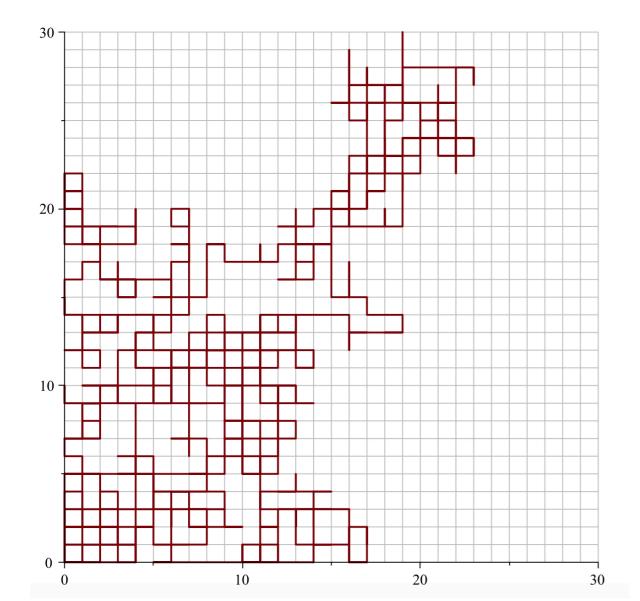


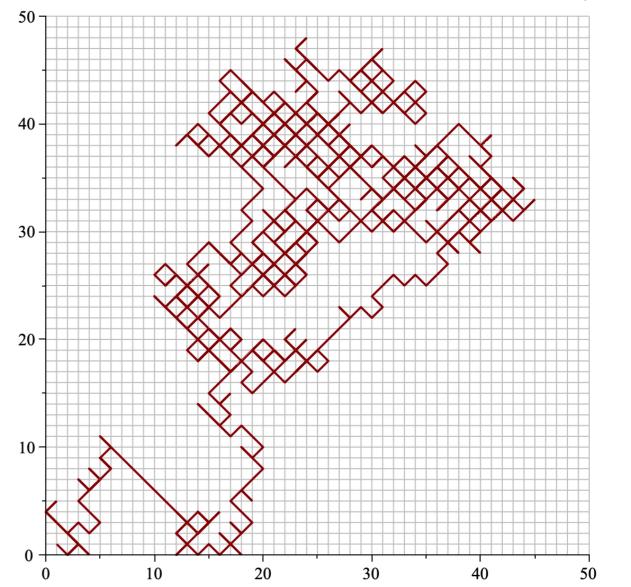
Walks in an Orthant

Uniform diagonal expression for walk models in \mathbb{N}^d whose step sets $\mathcal{S} \subset \{\pm 1, 0\}^d$ are symmetric over every axis.

walks
$$\sim |\mathcal{S}|^n \cdot n^{-d/2} \cdot \left(\left(s^{(1)} \cdots s^{(d)} \right)^{-1/2} \pi^{-d/2} |\mathcal{S}|^{d/2} + O\left(\frac{1}{n} \right) \right)$$

M. and Mishna, 2016





Lonesum Matrices (Khera, Lundberg, and M.)

A **lonesum matrix** is a 0-1 matrix that is uniquely determined by its row and column sums.

$$\begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \quad \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$$

$$\begin{pmatrix} 0 & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & 1 \end{pmatrix}$$
$$\begin{pmatrix} 1 & 1 \\ 0 & 0 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$$
$$\begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 0 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix} \begin{pmatrix} 1 & 1 \\ 1 & 1 \end{pmatrix}$$

NO

YES

Lonesum Matrices (Khera, Lundberg, and M.)

A **lonesum matrix** is a 0-1 matrix that is uniquely determined by its row and column sums.

$$F(x,y) = \sum_{n,k>0} \frac{B_{n,k}}{n!k!} x^n y^k = \frac{1}{e^{-x} + e^{-y} - 1}$$

Noncommutative Biology: Sequential Regulation of Complex Networks Letsou and Cai. PLOS Computational Biology, 2016.

Together with the fact that the reachable configurations are a subset of the staircase matrices, this implies that the reachable configurations and the lonesum matrices are in fact the same set, and we have

Theorem 3 The number of reachable configurations in the (n, m) ratchet network with $l_n = l_m = 1$ and threshold 1 scales as the poly-Bernoulli numbers $B_m^{-n} = B_n^{-m}$.

Lonesum Matrices (Khera, Lundberg, and M.)

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$$F(x,y) = \sum_{n,k>0} \frac{B_{n,k}}{n!k!} x^n y^k = \frac{1}{e^{-x} + e^{-y} - 1}$$

Let $f(t) = t/(1 - e^t) \log(1 - e^{-t})$.

Theorem. If $n, k \to \infty$ such that $n/k \to \lambda > 0$ then

$$B_{n,k} = \frac{a^{-n}b^{-n}}{\sqrt{k}} \frac{n!k!}{\sqrt{2\pi ae^{-a}[be^{-b} + ae^{-a} - ab]}} \left(1 + O\left(k^{-1}\right)\right),$$

where $a = f^{-1}(\lambda)$ and $b = f^{-1}(1/\lambda)$

Proof Idea of Main Theorem

Approach 1: Cones of Hyperbolicity (BP 2011)

- Criticality says something about the tangent space to $\mathcal V$
- Use this to locally deform around non-critical minimal points
- Glue these deformations together with roots of unity

Approach 2: Intersection Classes (BMP 2022)

- ullet Use geometric arguments to "flow" cycle on ${\mathcal V}$
- More on this approach tomorrow

ACSV Complexity Results

Suppose that $G(\mathbf{z})$ and $H(\mathbf{z})$ have coefficients $\leq 2^h$ and degree q Suppose also that the power series of $F(\mathbf{z})$ has non-negative coefficients

Theorem (M. and Salvy, 2016)

Under generic and verifiable assumptions one can **find all minimal critical points**, and compute asymptotics in $\tilde{O}(hq^{4d+5})$ bit operations.

Can remove non-negativity assumption, with increased complexity.

Theorem (M. and Salvy, 2021)

Under verifiable assumptions, one can find minimal critical points in $\tilde{O}(hq^{9d+4}2^{3d})$ bit operations.

ACSV Complexity Results

Theorem (M. and Salvy, 2021)

Under verifiable assumptions, one can find minimal critical points in $\tilde{O}(hq^{9d+4}2^{3d})$ bit operations.

General Idea:

- Assumptions imply finite number of critical points
- Use a univariate (Kronecker) representation to encode them
- Reduce everything to polynomial equalities and inequalities with bounded degrees and coefficient sizes
- Use numerical methods with sufficient precision to test minimality

Irrationality of Zeta(3)

Exercise

Be the first in your block to prove by a 2-line argument that $\zeta(3)$ is irrational.⁷

6 Given the definitions of 5 show that $a_n b_{n-1} - a_{n-1} b_n = b_n^{-3}$ and $b_n = O(\alpha^n)$ with $\alpha = (1 + \sqrt{2})^4$. Conclude that $\zeta(3)$ is irrational because $\log \alpha > 3$.

A Proof that Euler Missed, Alfred van der Poorten

$$b_n = \sum_{k=0}^{n} {n \choose k}^2 {n+k \choose k}^2 = [(xyzt)^n] \frac{1}{1 - t(1+x)(1+y)(1+z)(1+y+z+yz+xyz)}$$

Irrationality of Zeta(3)

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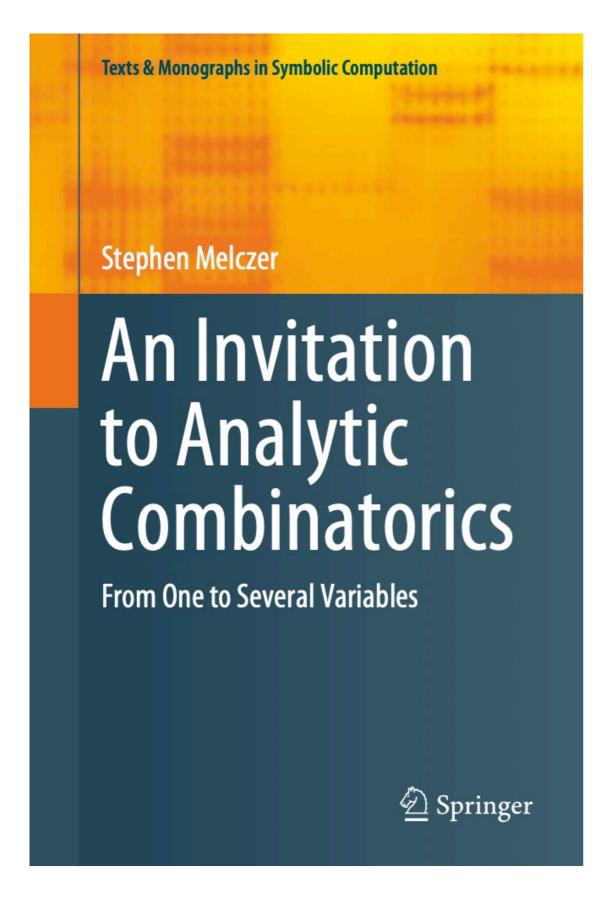
A, U, PRINT := DiagonalAsymptotics (numer (F), denom (F), [a,b,c,z],u,k, useFGb):
A, U;
$$\frac{1}{4} \frac{\left(\frac{2u-366}{34u+1458}\right)^k \sqrt{2} \sqrt{\frac{2u-366}{-96u-4192}}}{k^{3/2} \pi^{3/2}}, [RootOf(Z^2-366Z-17711, -43.27416997969]$$

Restricted Factors in Words

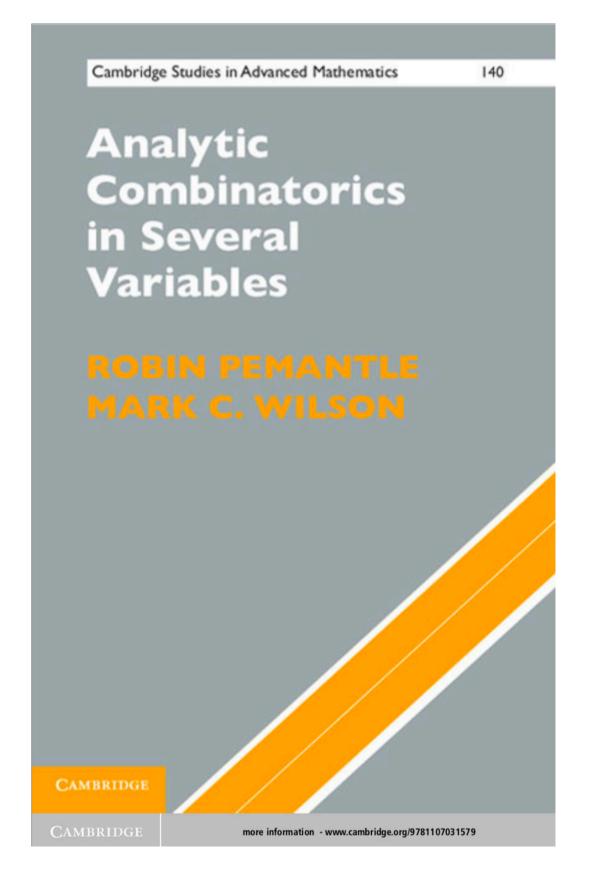
The number of balanced binary strings with no substring equal to 10101101 and 1110101 is the main diagonal of

$$\frac{1 - x^3y^6 + x^3y^4 + x^2y^4 + x^2y^3}{1 - x - y + x^2y^3 - x^3y^3 - x^4y^4 - x^3y^6 + x^4y^6}$$

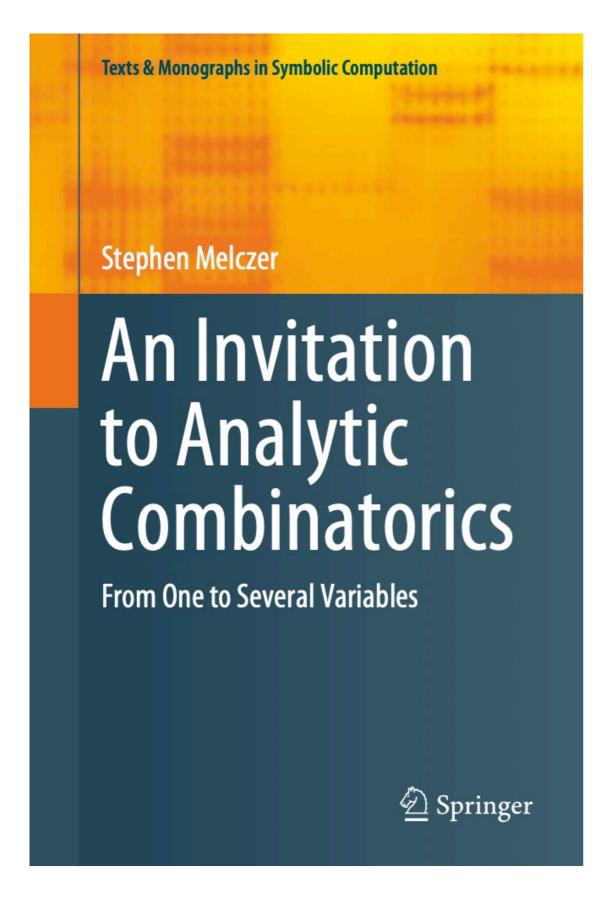
```
> ASM, U := Diagonal Asymptotics (numer (F), denom (F), indets (F), u, k, true, u-T, T): ASM; \frac{1}{2} \left( \frac{84 u^{20} + 240 u^{19} - 285 u^{18} - 1548 u^{17} - 2125 u^{16} - 1408 u^{15} + 255 u^{14} + 756 u^{13} + 2509 u^{12} + 2856 u^{11} + 605 u^{10} + 2020 u^9 + 1233 u^8 - 1760 u^7 + 212 u^{10} + 120 u^{19} + 258 u^{18} + 500 u^{17} + 440 u^{16} - 102 u^{15} - 378 u^{14} - 1544 u^{13} - 2142 u^{12} - 550 u^{11} - 2222 u^{10} - 1644 u^9 + 2860 u^8 - 1848 u^7 + 123 u^8 + 123 u^8
```



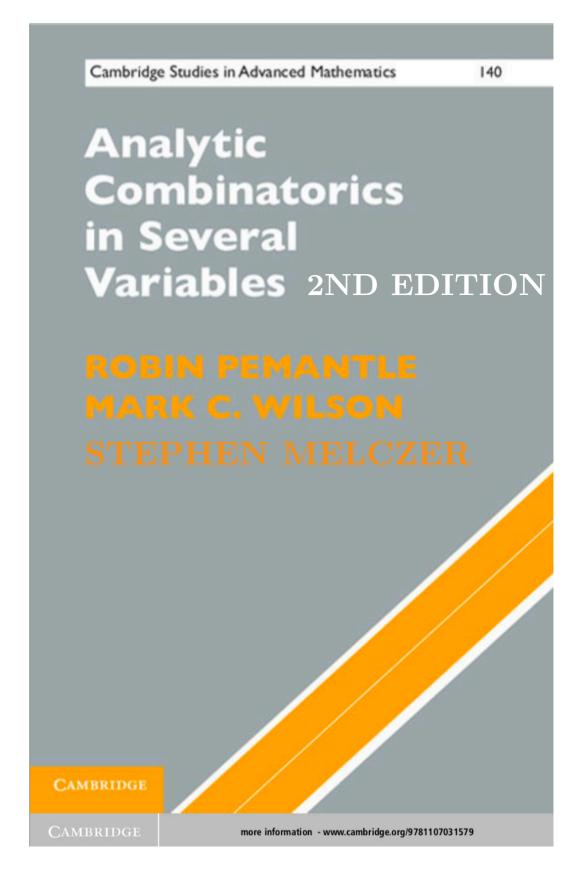
For new researchers — focus more on explicit results and computation



Most general theory, centred around topological framework



For new researchers — focus more on explicit results and computation



Most general theory, centred around topological framework

Implementations

- Sage package of Alex Raichev for computing asymptotic contributions of (already certified) minimal critical points
 - Package was not well maintained. Currently being fixed and extended to certify minimality by Hackl, Selover, and Wong.
- Julia implementation using Homotopy Continuation to certify minimality by Kisun Lee and Josip Smolić.
- Maple implementation of Melczer and Salvy for certifying minimality.
- Sage code developed for An Invitation to Analytic Combinatorics (available on melczer.ca/textbook)

Conclusion

Conclusion

- Analytic combinatorics is beautiful and powerful
- Analytic combinatorics in several variables is beautiful and powerful
- You don't need much more than univariate analytic combinatorics to get interesting results
- To remove some hard-to-check hypotheses we need to bring in new techniques and rely on some advanced mathematics

Lecture 2 (Tomorrow)

- Uniform asymptotics in varying directions
- Limit theorems
- (Some) non-smooth singular sets
- Morse-theoretic framework
- New applications
- And more!

THANK YOU!

An Invitation to Analytic Combinatorics melczer.ca/textbook

melczer.ca/ALEA22